

NVM - Refocusing on Memory Not Storage, Hierarchies

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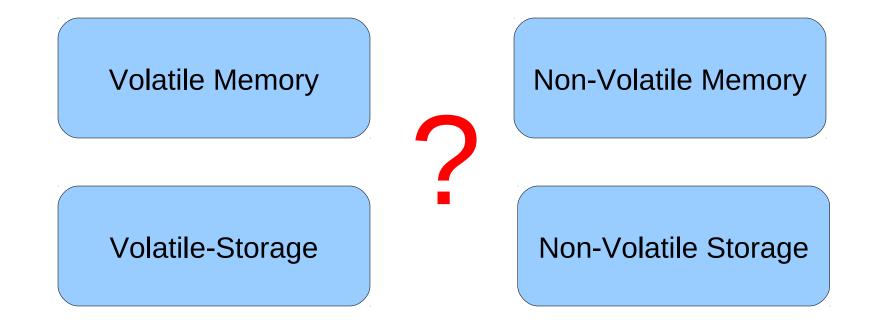
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Flash Memory Summit 2011 Santa Clara, CA



1









Von Neumann legacy means we have two classes

Volatile Memory

DRAM

Volatile-Storage

Time frame defines volatility

Flash Memory Summit 2011 Santa Clara, CA Non-Volatile Memory

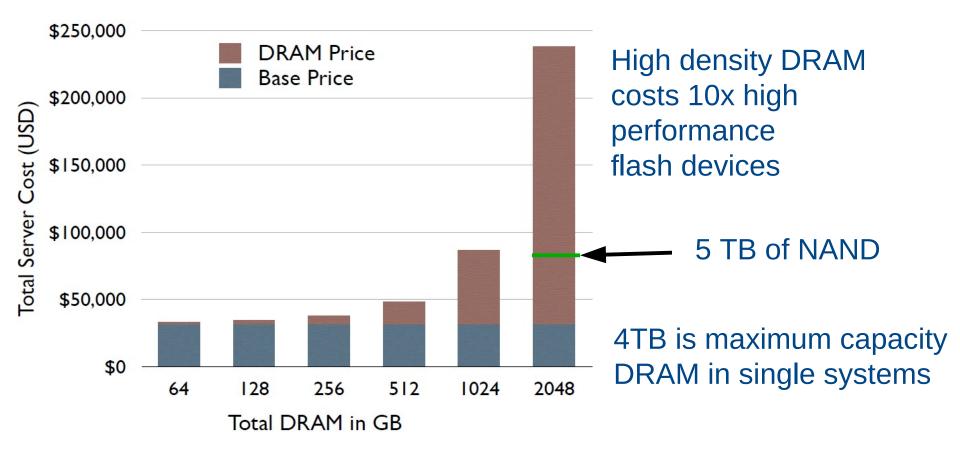
Hot OS research topic

Non-Volatile Storage

Magnetic Disks NAND based Disks Tape Drives Optical Media











High Density PCIe NAND-flash

5 rack units, 45TB capacity, 1.2kW power consumption

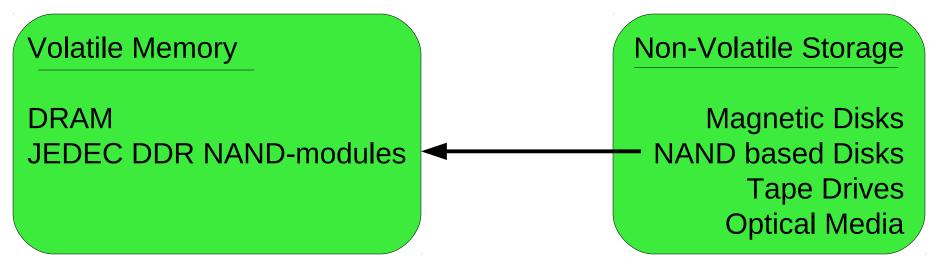
High Density DRAM

DRAM (GB)	128	256	512	1024	2048	4096
Space (RU)	6	6	10	40	80	80
Power (kW)	1.1	1.4	2.7	6.5	7.3	14.4





Approach 1: Move NAND-flash onto the memory bus

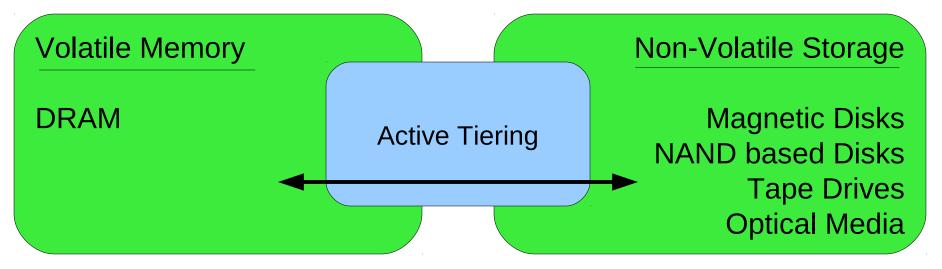


Advantages: Application Simplicity, Latency, Bandwidth Disadvantages: FTL Handling, Engineering Effort, Standards



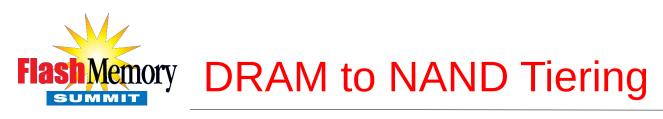


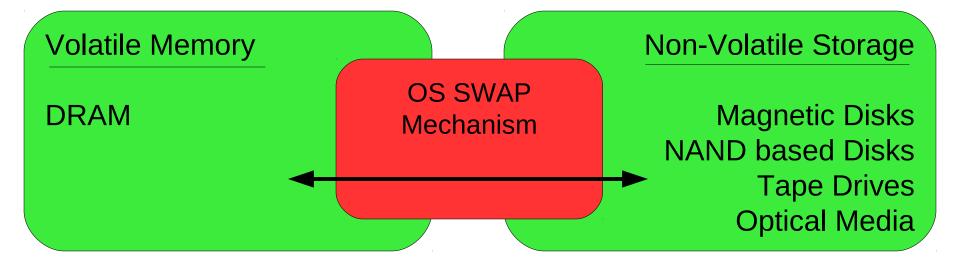
Approach 2: Allow tiering between DRAM and NAND



Advantages: Application Simplicity, No Device Engineering Leverage Faster Storage Advancements Disadvantages: Implementations Optimized for Magnetic Disk







Traditional SWAP: "Last resort" - before OOM <= 30MB/s throughput 10-100ms software overhead





Transparent Expansion of Application Memory *

Application Transparency: No source code modification!

Unhindered Access to DRAM

Low overhead tiering: Must not inhibit flash performance

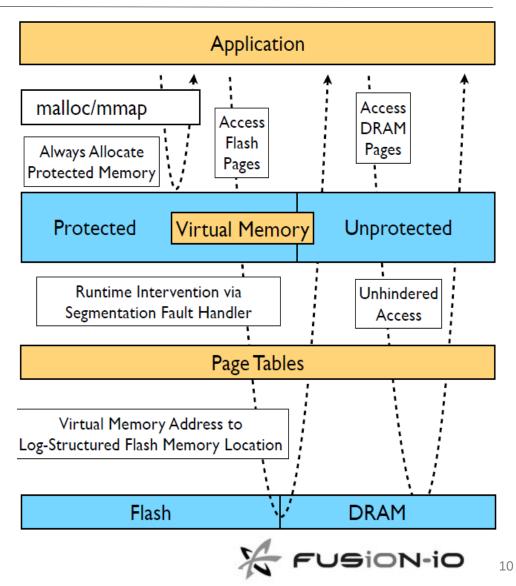
Intelligent paging decisions including application hints



DRAM and NAND Tiering

Builds on historical distributed shared memory concepts.

Instead of distributed DRAM, use locally or remote attached NAND-flash.





- 1. Allow thread-level paging: avoid process locking
- 2. Optimize for page-in operation to reduce latency
- 3. Intelligent utilization of flash-devices (wear-out aware)
- 4. Optional application hints to intelligently page data







Xeon 3.43Ghz with DDR3 1333Mhz running Linux

- 10,900,000 Random 64Byte Memory IOPS
- 120,000 Random 512B NAND-flash IOPS
- Linux SWAP: 11k Random NAND-memory IOPS
- TEAM Tiering: 93k Random NAND-memory IOPS

How does latency affect application performance?







Can we provide near-native application support?

Is transparency goal performance hindering?

Is NAND fast enough be a main memory replacement?

Is NAND fast enough to be used for tiered main memory?





Choose Percona MySQL 5.5 running TPC-C

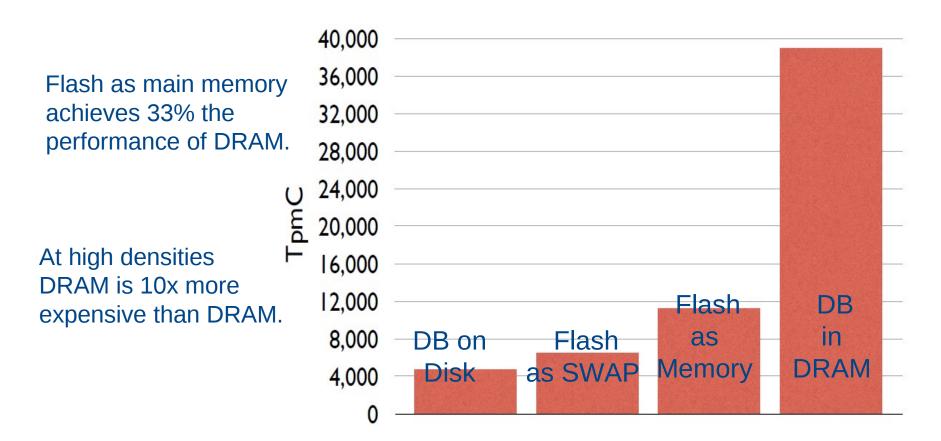
- Good native tiering support between DRAM and disk
- More memory is always better
- Not cost effective to put 100% of data on flash







Small Database Workloads

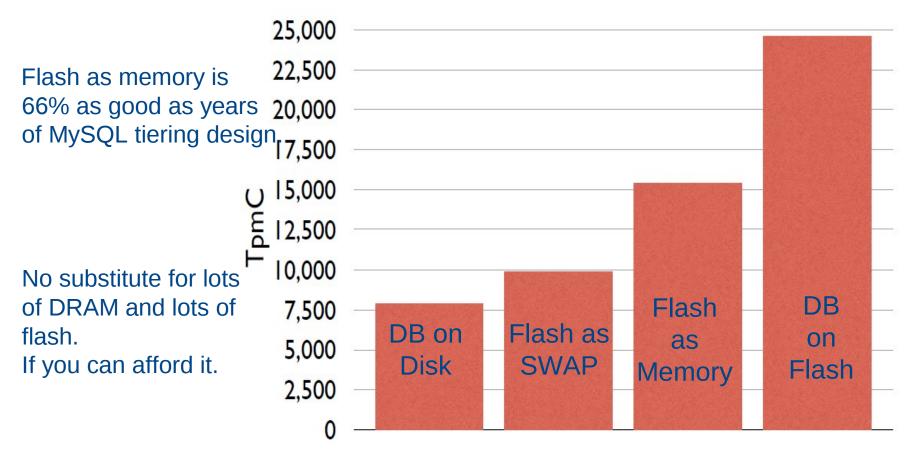


24 core Xeon, 40G DRAM, 140G Fusion-io NAND-flash: 40G DB size





Medium Database Workloads

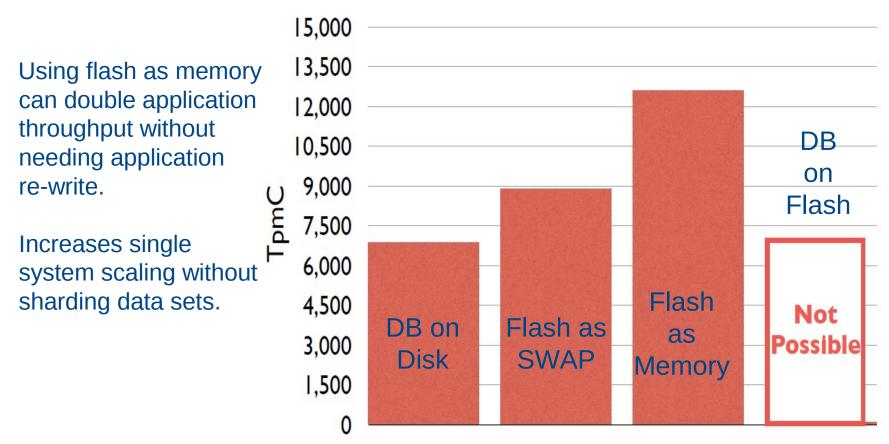


24 core Xeon, 40G DRAM, 140G Fusion-io NAND-flash: 140G DB size





Large Database Workloads



24 core Xeon, 40G DRAM, 140G Fusion-io NAND-flash: 400G DB size





Legacy applications may require 10's of TB of main memory. Scaling nodes up has been seen as non cost-effective. Continued sharding of data makes locality hard to maintain. A low-power, high density replacement for DRAM is needed.





NAND-flash is *not ready* as a wholesale DRAM replacment. Dense, power efficient, cheap. Too slow.

NAND-flash + DRAM tiering can provide:

66% the performance of an application re-write for tiering 33% the performance of all DRAM, 8% the TCO, and 5% power consumption.

NAND-flash is a *cost effective* way to build large memory systems.





Thank you!

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